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TECHNICAL SPECIFICATION

5G;
Specification of an example algorithm
for $f5^{}$ function for MILENAGE and Tuak;**
Algorithm specification and test data
(3GPP TS 35.249 version 19.1.0 Release 19)



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Foreword

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In the present document, modal verbs have the following meanings:

- shall** indicates a mandatory requirement to do something
- shall not** indicates an interdiction (prohibition) to do something

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- should** indicates a recommendation to do something
- should not** indicates a recommendation not to do something
- may** indicates permission to do something
- need not** indicates permission not to do something

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- can** indicates that something is possible
- cannot** indicates that something is impossible

The constructions "can" and "cannot" are not substitutes for "may" and "need not".

- will** indicates that something is certain or expected to happen as a result of action taken by an agency the behaviour of which is outside the scope of the present document
- will not** indicates that something is certain or expected not to happen as a result of action taken by an agency the behaviour of which is outside the scope of the present document
- might** indicates a likelihood that something will happen as a result of action taken by some agency the behaviour of which is outside the scope of the present document

might not indicates a likelihood that something will not happen as a result of action taken by some agency the behaviour of which is outside the scope of the present document

In addition:

is (or any other verb in the indicative mood) indicates a statement of fact

is not (or any other negative verb in the indicative mood) indicates a statement of fact

The constructions "is" and "is not" do not indicate requirements.

1 Scope

The present document presents an optional security enhancement for the MILENAGE [5, 6, 12, 13] and Tuak example algorithm set [7, 14, 15] for the 3GPP authentication and key agreement functions. The enhancement addresses a subscriber traceability issue discovered by academic researchers [9] and is provided as an alternative anonymity key generation function $f5^{**}$ that, when enabled, replaces the $f5^*$ function of the aforementioned example algorithm sets. A functionally equivalent security enhancement is already included in the new MILENAGE-256 example algorithm set [8].

2 References

The following documents contain provisions which, through reference in this text, constitute provisions of the present document.

- References are either specific (identified by date of publication, edition number, version number, etc.) or non-specific.
- For a specific reference, subsequent revisions do not apply.
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- [1] 3GPP TR 21.905: "Vocabulary for 3GPP Specifications".
- [2] 3GPP TS 33.102: "3rd Generation Partnership Project; Technical Specification Group Services and System Aspects; 3G Security; Security Architecture".
- [3] The Advanced Encryption Standard (AES), NIST FIPS 197, NIST, 2001.
- [4] Rijndael information page, NIST archived AES submissions, <https://csrc.nist.gov/projects/cryptographic-standards-and-guidelines/archived-crypto-projects/aes-development#rijndael>.
- [5] 3GPP TS 35.205: "Specification of the MILENAGE Algorithm Set: An example algorithm set for the 3GPP authentication and key generation functions $f1$, $f1^*$, $f2$, $f3$, $f4$, $f5$ and $f5^*$; Document 1: General".
- [6] 3GPP TS 35.206: "Specification of the MILENAGE Algorithm Set: An example algorithm set for the 3GPP authentication and key generation functions $f1$, $f1^*$, $f2$, $f3$, $f4$, $f5$ and $f5^*$; Document 2: Algorithm Specification".
- [7] 3GPP TS 35.231: "Specification of the TUAK algorithm set: A second example algorithm set for the 3GPP authentication and key generation functions $f1$, $f1^*$, $f2$, $f3$, $f4$, $f5$ and $f5^*$; Document 1: Algorithm specification".
- [8] ETSI SAGE, "Specification of the MILENAGE-256 algorithm set: An example set of 256-bit 3GPP Authentication and Key Generation functions $f1$, $f1^*$, $f2$, $f3$, $f4$, $f5$, $f5^*$ and $f5^{**}$, Document 1: General".
- [9] R. Borgaonkar, "New Privacy Threat on 3G, 4G, and Upcoming 5G AKA Protocols", in Proceedings on Privacy Enhancing Technologies 2019(3):108-127. Also available at <https://eprint.iacr.org/2018/1175.pdf> (published online: July 2019).
- [10] ETSI TR 133 909, "Report on the design and evaluation of the MILENAGE algorithm set; Deliverable 5: An example algorithm for the 3GPP authentication and key generation functions", 2001.
- [11] ETSI SAGE: "Specification of the MILENAGE-256 algorithm set: An example set of 256-bit 3GPP Authentication and Key Generation functions $f1$, $f1^*$, $f2$, $f3$, $f4$, $f5$, $f5^*$ and $f5^{**}$, Document 4: Summary and Results of Design and Evaluation".

- [12] 3GPP TS 35.207: "Specification of the MILENAGE Algorithm Set: An example algorithm set for the 3GPP authentication and key generation functions f_1 , f_1^* , f_2 , f_3 , f_4 , f_5 and f_5^* ; Document 3: Implementors' test data".
- [13] 3GPP TS 35.208: "Specification of the MILENAGE Algorithm Set: An example algorithm set for the 3GPP authentication and key generation functions f_1 , f_1^* , f_2 , f_3 , f_4 , f_5 and f_5^* ; Document 4: Design conformance test data".
- [14] 3GPP TS 35.232: "Specification of the TUAK algorithm set: A second example algorithm set for the 3GPP authentication and key generation functions f_1 , f_1^* , f_2 , f_3 , f_4 , f_5 and f_5^* ; Document 2: Implementors' test data".
- [15] 3GPP TS 35.233: "Specification of the TUAK algorithm set: A second example algorithm set for the 3GPP authentication and key generation functions f_1 , f_1^* , f_2 , f_3 , f_4 , f_5 and f_5^* ; Document 3: Design conformance test data".

3 Definitions of terms, symbols and abbreviations

3.1 Terms

For the purposes of the present document, the terms given in TR 21.905 [1] and the following apply. A term defined in the present document takes precedence over the definition of the same term, if any, in TR 21.905 [1].

3.2 Symbols

For the purposes of the present document, the following symbols apply:

\oplus	The bitwise exclusive-OR operation
$X[i]$	The i^{th} bit of the variable X . ($X = X[0] \parallel X[1] \parallel X[2] \parallel \dots$).
\parallel	Concatenation: $X[0]X[1] \dots X[n-1] \parallel Y[0]Y[1] \dots Y[m-1] = X[0]X[1] \dots X[n-1]Y[0]Y[1] \dots Y[m-1]$
$\text{rot}(x,r)$	The result of cyclically rotating the 128-bit value x by r bit positions towards the most significant bit. See [6].
MILENAGE-128	The MILENAGE algorithm set as defined in [5, 6]. NOTE: The term MILENAGE-128 is used to distinguish this algorithm set from the more recently defined MILENAGE-256 set [8].
$E[x]_k$	The AES-based cryptographic kernel of MILENAGE-128. Specifically, the result of applying the block cipher E (AES) to the input value x using the key k .
Tuak	The Tuak algorithm set as defined in [7].
Π	The Keccak-based cryptographic kernel of Tuak.

3.3 Abbreviations

For the purposes of the present document, the abbreviations given in TR 21.905 [1] and the following apply. An abbreviation defined in the present document takes precedence over the definition of the same abbreviation, if any, in TR 21.905 [1].

AES	Advanced Encryption Standard
AKA	Authentication and Key Agreement
MAC	Message Authentication Code
ETSI SAGE	ETSI Security Algorithms Group of Experts

4 Structure of this specification

This specification is organised as follows:

Clause 3 and 5 introduce symbols and notation used in the subsequent clauses.

Clause 6 provides a summary of all variables (inputs, outputs, and intermediary values) used in the algorithm specification.

Clause 7 provides a general definition of the $f5^{**}$ function which is in agreement with the API of [2, Annex J] and which is independent on the underlying cryptographic kernel, i.e. it applies to both the instantiation within the MILENAGE-128 framework as well as with Tuak.

Clause 8 and 9 contains the specifics of the MILENAGE-128 and Tuak instantiations, respectively.

Clause 10 provides a security analysis.

Clauses 11 contains test data for implementors of the MILENAGE-128 and Tuak-based variants.

NOTE: The reader is reminded of the fact that both MILENAGE-128, as well as Tuak, in principle allow distinct choices for the security primitive (block cipher or hash function, respectively) employed as cryptographic kernel, provided the same input/output parameters can be securely supported. However, the present document assumes that the default kernels are used, i.e. that AES/Rijndael [3, 4] is used to instantiate MILENAGE-128 [5, 6] and the Keccak-based kernel of [7] is used to instantiate Tuak.

5 Introductory information

5.1 Notation

5.1.1 Radix

Unless otherwise noted, integer values are represented in decimal. The prefix **0x** indicates **hexadecimal** integers.

5.1.2 Bit ordering and arrays

The same ordering as in [5, 6, 7] shall apply. In other words, all data variables in the present document are presented with the most significant substring on the left hand side and the least significant substring on the right hand side. A substring may be a bit, byte or other arbitrary length bit string. Where a variable is broken down into a number of substrings, the left-most (most significant) substring is numbered 0, the next most significant is numbered 1, and so on through to the least significant.

EXAMPLE: **RAND**[0] is the most-significant bit of **RAND** and **RAND**[127] is the least significant bit of **RAND**.

6 List of variables

6.1 General

The set of variables defined in [5, 6, 7] shall apply also for the purposes of the present document. The following variables are highlighted since they are of particular relevance for the definition of $f5^{**}$.

Table 6.1-1: List of variables

Name	Comment
AK	An anonymity key that is output by the functions f_5 and f_5^* of [5, 6]. NOTE: The MILENAGE-256 algorithm set [8] distinguishes the AK output by f_5 from those output by f_5^* by using the variable AK* for the latter function. This distinction is not made in [5, 6, 7] and is therefore not made in the present document either, to avoid confusion when cross-referencing these documents.
AMF	A two byte of authentication management field that is input to the functions f_1 and f_1^* .
K	A subscriber key that is an input to the functions f_1 , f_1^* , f_2 , f_3 , f_4 , f_5 , f_5^* .
MAC-S	A resynchronisation authentication code that is output by the function f_1^* .
RAND	A random challenge that is an input to the functions f_1 , f_1^* , f_2 , f_3 , f_4 , f_5 , f_5^* .
SQN	A sequence number that is an input to either of the functions f_1 and f_1^* . (For f_1^* this input is more precisely called SQN_{MS})
c1, ..., c5 and r1, ..., r5	Operator selectable, 128-bit algorithm customization constants used in the definition of MILENAGE-128. These are not used in Tuak.
TEMP	An intermediate 128-bit value occurring during MILENAGE-128 computation. This value is not used by Tuak.
IN	Input value formed during the computation of the f-functions. IN is 128 bits in MILENAGE-128 and 1600 bits in Tuak. NOTE: Two specific IN -values of importance for the MILENAGE-128 variant of f_5^{**} are below denoted IN1 and IN6 .
OUT	Output value obtained by invoking the cryptographic kernels of MILENAGE-128 or Tuak on the IN values, and from which the f-function outputs are obtained by selecting some or all bits.
INSTANCE	An 8-bit value used as part of the input values (IN) to Tuak and which is assigned different values for the different f-functions. This value is not used by MILENAGE-128.
OPc	An algorithm customization value (derived from an operator selected value OP).

6.2 Input/output variables specific for f_5^{**}

Table 6.2-1: Input/output variables specific for f_5^{**}

Name	Comment
AK	This variable name is reused to denote the anonymity resynchronisation key that is output by the function f_5^{**} . NOTE: Under normal operation, the AK value is provided by f_5 . If a resynchronisation procedure is required, an additional AK value is provided either by f_5^* , or by f_5^{**} .
K	The subscriber key is an input also to f_5^{**} .
MAC-S	The resynchronisation authentication code that is output by the function f_1^* and which is used also as an input to f_5^{**} .
RAND	The random challenge, also used as an input to f_5^{**} .

6.3 Input/output parameter sizes for $f5^{**}$

Table 6.3-1: Input/output parameter sizes for $f5^{**}$

Name	Permitted values	Comment
AK_{sz}	48	The size in bits of the anonymity key AK .
K_{sz}	Kernel-dependent	The size in bits of the subscriber key K .
MAC_{sz}	Kernel-dependent	The size in bits of the resynchronisation authentication code MAC-S .
$RAND_{sz}$	128	The size in bits of the random challenge RAND .
SQN_{sz}	48	The size in bits of the sequence number SN .

7 General definition of $f5^{**}$

Annex J of [2] specifies an alternative $f5^*$ function offering protection against the resynchronization attack of [9]. Its API is specified as follows:

Inputs: the subscriber key **K**, the network challenge **RAND**, and the resynchronisation MAC-value **MAC-S**.

Output: an anonymity key, **AK**.

In the present document, this function is defined as:

$AK = f5^{**}(K, RAND, MAC-S)$.

NOTE: $f5^{**}$ can also depend on other values, e.g. operator selected customization values, not shown.

8 MILENAGE-128 based $f5^{**}$

8.1 Supported parameter sizes

The MILENAGE-128 version of $f5^{**}$ shall support $K_{sz} = 128$ bits and $MAC_{sz} = 64$ bits only.

8.2 Algorithm specification

8.2.1 General

The choice to implement $f5^{**}$ instead of $f5^*$ has no impact on the f-functions ($f1$, $f1^*$, $f2$, $f3$, $f4$, and $f5$). The specifications [5, 6] shall apply except as specified in clause 8.2.2 below.

8.2.2 $f5^{**}$ specification

Precisely the following modifications to clause 4.1 of [6] shall apply:

- References to the function $f5^*$ shall be ignored.
- An additional 128-bit value **IN6** shall be computed as follows:

$$IN6[0] IN6[1] \dots IN6[63] = (MAC-S[0] \oplus 1) (MAC-S[1] \oplus 1) MAC-S[2] \dots MAC-S[63]$$

and

$$IN6[64] IN6[65] \dots IN6[127] = MAC-S[0] MAC-S[1] \dots MAC-S[62] MAC-S[63].$$

- An additional value **OUT6** shall be computed as follows:

$$OUT6 = E[TEMP \oplus \text{rot}(IN6 \oplus OP_C, r1) \oplus c1]_K \oplus OP_C.$$

4. Output of $f5^{**} = \mathbf{AK}$, where $\mathbf{AK}[0] .. \mathbf{AK}[47] = \mathbf{OUT6}[0] .. \mathbf{OUT6}[47]$

NOTE 1: By modification 2, the two 64-bit halves of $\mathbf{IN6}$ are identical, except that the two bits $\mathbf{IN6}[0]$, $\mathbf{IN6}[1]$ are both inverted, relative to $\mathbf{IN6}[64]$, $\mathbf{IN6}[65]$.

NOTE 2: Conceptually, modification 3 defines two new constants $\mathbf{c6}$ and $\mathbf{r6}$, but with values identical to $\mathbf{c1}$ and $\mathbf{r1}$.

9 Tuak based $f5^{**}$

9.1 Supported parameter sizes

The Tuak version of $f5^{**}$ shall support $K_{SZ} = 128$ and 256 bits and shall support $MAC_{SZ} = 64, 128, \text{ and } 256$ bits. The chosen values of these quantities shall be the same as those commonly chosen for the complete algorithm set.

9.2 Algorithm specification

9.2.1 General

The choice to implement $f5^{**}$ instead of $f5^*$ has no impact on the f -functions $f1, f1^*, f2, f3, f4, \text{ and } f5$. The specification [7] shall apply except as specified in clause 7.2.2 below.

9.2.2 $f5^{**}$ specification

Clause 6.5 of [7] shall be replaced by the following. Recall that the bit-length of $\mathbf{MAC-S}$ is defined as MAC_{SZ} .

The internal **INSTANCE** variable shall be constructed as follows:

$\mathbf{INSTANCE}[0] .. \mathbf{INSTANCE}[1] = 1,1$

$\mathbf{INSTANCE}[2] .. \mathbf{INSTANCE}[4] = 0,0,1$ if MAC_{SZ} is 64 bits,

$= 0,1,0$ if MAC_{SZ} is 128 bits,

$= 1,0,0$ if MAC_{SZ} is 256 bits

$\mathbf{INSTANCE}[5] .. \mathbf{INSTANCE}[6] = 0,0$

$\mathbf{INSTANCE}[7] = 0$ if the length of \mathbf{K} is 128 bits

$= 1$ if the length of \mathbf{K} is 256 bits.

Given **INSTANCE** as above, this value shall be assigned to $\mathbf{IN}[256] .. \mathbf{IN}[263]$. The rest of the 1600-bit value \mathbf{IN} (i.e. $\mathbf{IN}[0].. \mathbf{IN}[255]$ and $\mathbf{IN}[264].. \mathbf{IN}[1599]$) shall then be constructed in exactly the same way as for $f2, f3, f4, \text{ and } f5$ except that

$\mathbf{IN}[768] .. \mathbf{IN}[768 + MAC_{SZ} - 1] = \mathbf{MAC-S}[MAC_{SZ} - 1] .. \mathbf{MAC-S}[0]$,

$\mathbf{IN}[i] = 0$, for $768 + MAC_{SZ} \leq i \leq 1023$,

$\mathbf{IN}[i] = 1$ for $1024 \leq i \leq 1028$.

Then the permutation:

$\mathbf{OUT} = \Pi(\mathbf{IN})$

is applied and the function output is extracted as follows:

Output of $f5^{**} = \mathbf{AK}$, where:

$\mathbf{AK}[0] .. \mathbf{AK}[47] = \mathbf{OUT}[815] .. \mathbf{OUT}[768]$.

NOTE 1: The value **INSTANCE**[0] **INSTANCE**[1] is identical to that defined for f5* but **INSTANCE**[2] .. **INSTANCE**[4] is different.

NOTE 2: For f2 to f5, the input padding starts at bit **IN**[768], while for f5**, the padding starts after **MAC-S** has been zero-padded to 256 bits, i.e. in bit **IN**[1024].

10 Security analysis

10.1 Resynchronisation attack description and impact

The attack presented in [8] can be summarised as follows. An attacker, Eve, has previously recorded an authentication between a victim, Bob, and the network. This resulted in a resynchronisation procedure. (In fact, Eve can force such a resynchronisation by replaying a previously observed authentication challenge, known to be directed to Bob.) Thus, Eve has at time T1 obtained a value of the form

$$\mathbf{SQN}_{\text{Bob}}^{(T1)} \oplus f5^*(\mathbf{K}_{\text{Bob}}, \mathbf{RAND}^{(T1)})$$

as well as the network authentication token **AUTN**^(T1), used at that time.

Later, at time T2, Eve wants to test if a certain subscriber is Bob or someone else, so she replays (**RAND**^(T1), **AUTN**^(T1)).

If it is not Bob, the UE/USIM will generate a "network authentication failure", observable to Eve and she can exclude that it is Bob.

On the other hand, if it is Bob, Bob's UE/USIM will initiate a resynchronisation (also observable to Eve) and this will reveal $\mathbf{SQN}_{\text{Bob}}^{(T2)} \oplus f5^*(\mathbf{K}_{\text{Bob}}, \mathbf{RAND}^{(T1)})$. Taking the XOR of these values, Eve obtains

$$\mathbf{SQN}_{\text{Bob}}^{(T1)} \oplus \mathbf{SQN}_{\text{Bob}}^{(T2)}.$$

Assuming Eve knows the sequence number generation method, at least some bits of this value will have a low Hamming weight since the values are likely related. In fact, if Bob has not performed any authentication since time T1, the XOR-value will be identically zero.

In conclusion, this provides Eve with a tracing mechanism for certain targeted subscribers.

10.2 Analysis

10.2.1 General

To mitigate the attack, it is possible to use a more sophisticated concealment mechanism for the **SQN** value than a simple XOR. However, this requires quite substantial modifications.

Another possibility is to add some freshness to the input of f5*, making it likely that the outputs are different. Without major changes to the AKA protocol, the only source of such freshness is the **MAC-S** value, which has led to an alternative f5* function of the form

$$\mathbf{AK} = f5^*(\mathbf{K}, \mathbf{RAND}, \mathbf{MAC-S})$$

akin to the two variants defined in the present document. There still remain two "bad" events that would allow Eve to trace Bob:

1. The **MAC-S** values are derived by f1*, with dependence on **RAND** and **SQN**, and could still collide "by chance" even if the (**RAN**, **SQN**)-values are distinct. This however has a low probability, exponentially small in the bit-size of **MAC-S**.
2. As noted above, if the resynchronisation occurs without Bob having performed any successful authentication in the time-interval (T1, T2), **SQN** will still have the same value, causing **AK**-values to collide, which Eve can detect. There seems however to be no way to avoid this without substantial changes to the AKA protocol, and

therefore the $f5^{**}$ approach above has been deemed a good compromise with sufficient protection and low additional complexity.

NOTE: This approach to the resynchronisation protection is conceptually very similar to known solutions that protect stream-ciphers against nonce-reuse and the construction is therefore judged to be cryptographically sound. More discussion on this can be found in [11].

10.2.2 MILENAGE-128 based $f5^{**}$

The main issue is the analysis of the possibility that $f5^{**}$ output could collide with one of the other f-functions. This is in turn dependent on the distinctness of the $f5^{**}$ input:

$$\text{TEMP} \oplus \text{rot}(\mathbf{IN6} \oplus \text{OPC}, r1) \oplus c1 \quad (1)$$

from values of form

$$\text{TEMP} \oplus \text{rot}(\mathbf{IN1} \oplus \text{OPC}, r1) \oplus c1, \quad (2)$$

i.e. the input to $f1/f1^*$, as well the distinctness between (1) and values of form

$$\text{rot}(\text{TEMP}, r_i) \oplus c_i, \quad i = 2, 3, 4, 5, \quad (3)$$

i.e. the inputs occurring when computing $f2$ - $f5$. This is analysed in great detail in [10], an analysis which largely applies also here. Let 3^{128} be the 128-bit representation of the integer value $3 * 2^{126}$ i.e. the binary string 1100...0.

- The distinctness of (1) relative to the values (3) is ensured by parity requirements:
 - o Rotation does not affect parity.
 - o As required by [6], $c2, c3, c4, c5$ all have odd parity, while $c1$ has even parity.
 - o The value $\mathbf{IN6}$ has structure $(Y \parallel Y) \oplus 3^{128}$ (where $Y = \mathbf{MAC-S}$). The part $(Y \parallel Y)$ always has even parity and this is preserved by the exclusive-or with 3^{128} .
 - o Together, this implies that regardless of the parity of TEMP , the values involved in (1) and (3) can never have equal parity and are therefore distinct.
- The distinctness from the value $\text{TEMP} \oplus \text{rot}(\mathbf{IN1} \oplus \text{OPC}, r1) \oplus c1$ of (2) follows since $\mathbf{IN1}$ by [6] has format $\text{AMF} \parallel \text{SQN} \parallel \text{AMF} \parallel \text{SQN}$, i.e. is a value of format $X \parallel X$. But such a value can never be identical to a value of format $(Y \parallel Y) \oplus 3^{128}$, such as $\mathbf{IN6}$. Since the same $(c1, r1)$ -value is used both for $f5^{**}$ and $f1/f1^*$, the distinctness follows.

In conclusion, collision probabilities are not affected by $f5^{**}$.

10.2.3 Tuak based $f5^{**}$

Also here, the only additional concern is the possibility that the input to $f5^{**}$ could collide with the input when computing one of the other f-functions. The absence of such collisions is ensured by the **INSTANCE**-values included in the inputs, which are always distinct.

While $f5^{**}$ uses the same value for **INSTANCE**[0]**INSTANCE**[1] as $f5^*$ does, the values of **INSTANCE**[2]..**INSTANCE**[4] are never identical for $f5^*$ and $f5^{**}$ (they are always identically zero for $f5^*$).

NOTE: Furthermore, only one of $f5^*$ and $f5^{**}$ is assumed to be supported.

11 Implementors' test data

11.1 MILENAGE-128 variant

TEST SET #1

K 465b5ce8 b199b49f aa5f0a2e e238a6bc
 RAND 23553cbe 9637a89d 218ae64d ae47bf35
 SQN ff9bb4d0 b607
 AMF b9b9
 OP cdc202d5 123e20f6 2b6d676a c72cb318
 OPc cd63cb71 954a9f4e 48a5994e 37a02baf
 f1 4a9ffac3 54dfafb3
 f1* 01cfaf9e c4e871e9
 f2 a54211d5 e3ba50bf
 f5 aa689c64 8370
 f3 b40ba9a3 c58b2a05 bbf0d987 b21bf8cb
 f4 f769bcd7 51044604 12767271 1c6d3441
 f5* 451e8bec a43b
 f5** 4edd7fbd c382

TEST SET #2

K 0396eb31 7b6d1c36 f19c1c84 cd6ffd16
 RAND c00d6031 03dcee52 c4478119 494202e8
 SQN fd8eef40 df7d
 AMF af17
 OP ff53bade 17df5d4e 793073ce 9d7579fa
 OPc 53c15671 c60a4b73 1c55b4a4 41c0bde2
 f1 5df5b318 07e258b0
 f1* a8c016e5 1ef4a343
 f2 d3a628ed 988620f0
 f5 c4778399 5f72
 f3 58c433ff 7a7082ac d424220f 2b67c556
 f4 21a8c1f9 29702adb 3e738488 b9f5c5da
 f5* 30f11970 61c1
 f5** 7b958d44 d816

TEST SET #3

K fec86ba6 eb707ed0 8905757b 1bb44b8f
 RAND 9f7c8d02 1accf4db 213ccff0 c7f71a6a
 SQN 9d027759 5ffc
 AMF 725c
 OP dbc59adc b6f9a0ef 735477b7 fadf8374
 OPc 1006020f 0a478bf6 b699f15c 062e42b3
 f1 9cabc3e9 9baf7281
 f1* 95814ba2 b3044324
 f2 8011c48c 0c214ed2
 f5 33484dc2 136b
 f3 5dbdbb29 54e8f3cd e665b046 179a5098
 f4 59a92d3b 476a0443 487055cf 88b2307b
 f5* deacdd84 8cc6
 f5** e7fd8260 d2c9

TEST SET #4

K 9e5944ae a94b8116 5c82fbf9 f32db751
 RAND ce83dbc5 4ac0274a 157c17f8 0d017bd6
 SQN 0b604a81 eca8
 AMF 9e09
 OP 223014c5 806694c0 07ca1eee f57f004f
 OPc a64a507a e1a2a98b b88eb421 0135dc87
 f1 74a58220 cba84c49
 f1* ac2cc74a 96871837
 f2 f365cd68 3cd92e96
 f5 f0b9c08a d02e
 f3 e203edb3 971574f5 a94b0d61 b816345d
 f4 0c4524ad eac041c4 dd830d20 854fc46b
 f5* 6085a86c 6f63
 f5** f0e42c7f af58

TEST SET #5

K 4abldeb0 5ca6ceb0 51fc98e7 7d026a84
 RAND 74b0cd60 31a1c833 9b2b6ce2 b8c4a186
 SQN e880alb5 80b6
 AMF 9f07
 OP 2d16c5cd 1fdf6b22 383584e3 bef2a8d8
 OPc dcf07cbd 51855290 b92a07a9 891e523e

```
f1 49e785dd 12626ef2
f1* 9e857903 36bb3fa2
f2 5860fc1b ce351e7e
f5 31e11a60 9118
f3 7657766b 373d1c21 38f307e3 de9242f9
f4 1c42e960 d89b8fa9 9f2744e0 708ccb53
f5* fe2555e5 4aa9
f5** 7468dacf 4f72
```

```
TEST SET #6
K 6c38a116 ac280c45 4f59332e e35c8c4f
RAND ee6466bc 96202c5a 557abbef f8babf63
SQN 414b9822 2181
AMF 4464
OP 1ba00a1a 7c6700ac 8c3ff3e9 6ad08725
OPc 3803ef53 63b947c6 aaa225e5 8fae3934
f1 078adfb4 88241a57
f1* 80246b8d 0186bcf1
f2 16c8233f 05a0ac28
f5 45b0f69a b06c
f3 3f8c7587 fe8e4b23 3af676ae de30ba3b
f4 a7466cc1 e6b2a133 7d49d3b6 6e95d7b4
f5* 1f53cd2b 1113
f5** 74f55f33 b347
```

11.2 Tuak variant

Inputs to these test data are identical to those 6 test sets given in Clauses 6.3 to 6.8 of TS 35.232 [14]. Here only the delta-additions related to f5** function are listed.

```
=====
6.3 Test set 1 [3GPP TS 35.232 V17.0.0 (2022-03)]
=====
+Intermediate Values:
IN when computing f5**:
ff cb cc 40 1f 5d b4 3e a7 05 53 11 0c 33 e2 a8 23 46 95 ad c2 7a 83 5d 3c 51 87 0e
53 d9 04 bd c8 30 2e 31 4b 41 55 54 42 42 42 42 42 42 42 42 42 42 42 42 42 42 42 42
00 00 00 00 00 00 00 00 00 ab ab ab ab ab ab ab ab ab ab ab ab ab ab ab ab 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00
OUT when computing f5**:
9e 37 8b 43 2b 04 6f 39 40 f4 d9 e4 32 27 0d da 10 bc 5c 9f d6 c7 fb f7 5a 44 7f e2
16 a2 23 e6 d4 08 3f 37 84 e7 1f 28 84 77 7c d8 94 b7 71 9d 5d b8 db 1e 19 a0 20 2d
2f a4 ab 1b 24 8c 89 9e fe 23 d2 c8 31 e5 1e c4 80 58 8a ba 7c b8 80 c7 ce 2f 15 20
8a 91 26 be 7e 5f a5 d1 8e 5d d1 82 4f 66 18 a4 62 7d 0c 1f 72 37 8d 45 35 d7 26 25
13 6f ef 93 0a a5 45 f9 71 ff 5a e3 08 98 d5 a6 bf a2 97 e9 a1 b9 92 a6 9d d8 ad 80
e0 2b d0 1a 00 9a 31 79 8e eb 4f fd 94 15 00 65 99 86 d8 ae 44 6c d0 e4 7c 9d 45 ac
30 4a f2 5e e5 49 61 f2 3e 1d 18 59 60 b4 b2 eb dc 77 0f 1a 35 38 be e7 53 ec c5 7b
2b 06 73 f8

+Output Parameters:
f5**: 7d62a418664f
```

=====
6.4 Test set 2 [3GPP TS 35.232 V17.0.0 (2022-03)]
=====

+Intermediate Values:

IN when computing f5**:

24 85 7f 97 d5 74 59 de 65 0c 9d b2 c6 c6 36 0c 5d c9 72 ea 94 b2 a4 c8 03 c5 18 7e
42 25 54 30 d1 30 2e 31 4b 41 55 54 ef cd ab 89 67 45 23 01 ef cd ab 89 67 45 23 01
00 00 00 00 00 00 00 00 e0 e1 e2 e3 e4 e5 e6 e7 e8 e9 ea eb ec ed ee ef f0 f1 f2 f3
f4 f5 f6 f7 f8 f9 fa fb fc fd fe ff 5b 74 a0 7f 53 fa ea 6c 2c 84 f7 90 72 af 81 ef
00
00
00
00 00 00 00

OUT when computing f5**:

c6 58 e8 84 cb bf 0d 05 b8 26 ae a7 ff e6 5b 48 b9 f2 70 b1 39 ae 2b d0 79 94 93 bc
a4 9f 54 95 23 b8 87 4c 6a 72 d1 87 ae 44 0d 16 6a f8 31 ed 7c 81 f7 4a 00 03 1e fd
4a 86 a7 90 a0 ba 05 10 90 af 78 a1 2d 25 97 3b 77 00 44 17 8c 6e 52 d2 16 fd 71 08
6e 56 98 fe 62 7b ce 11 98 eb 05 94 83 07 69 7a 46 be 7a 25 52 c9 00 2c 6c 62 9d ab
81 da 76 4f 7f dd 16 c4 c0 8e fb 09 d8 bd 12 77 a5 b6 68 4b 95 92 4c e6 7f 88 38 bf
a4 91 a6 e0 ef 6a 92 eb c0 63 74 96 24 d3 ae 9c 91 06 d6 b8 33 76 3e 4a 16 0c 9f 28
1b 98 63 ce c3 97 e9 9a ed 40 49 77 cc 05 c2 4a 18 b4 0e d3 52 d5 92 97 07 b7 f0 89
66 33 04 97

+Output Parameters:

f5** : be467a690783

=====
6.5 Test set 3 [3GPP TS 35.232 V17.0.0 (2022-03)]
=====

+Intermediate Values:

IN when computing f5**:

24 85 7f 97 d5 74 59 de 65 0c 9d b2 c6 c6 36 0c 5d c9 72 ea 94 b2 a4 c8 03 c5 18 7e
42 25 54 30 e1 30 2e 31 4b 41 55 54 ef cd ab 89 67 45 23 01 ef cd ab 89 67 45 23 01
00 00 00 00 00 00 00 00 e0 e1 e2 e3 e4 e5 e6 e7 e8 e9 ea eb ec ed ee ef f0 f1 f2 f3
f4 f5 f6 f7 f8 f9 fa fb fc fd fe ff 21 b8 11 81 25 50 15 5b 23 fe c9 64 a1 d4 6f 9a
90 f3 99 64 21 c5 f8 54 6c a8 e3 c6 07 bf 7b 42 1f 00 00 00 00 00 00 00 00 00 00 00
00
00
00 00 00 00

OUT when computing f5**:

d9 17 a7 73 ld d7 80 7c db 9e 36 ba b1 70 34 a3 b6 1c 4f 85 e1 c2 ac 37 dd bb f4 43
e4 1a 69 a7 f4 a7 e1 a3 22 20 90 20 ef 80 d0 c4 06 34 40 f3 11 b9 29 41 fd 03 72 12
da 5d be d0 31 1a f8 c6 1c 83 f4 28 c1 0f ec 69 fd 46 d9 54 a9 64 6a d9 7b 35 fe 82
2c 35 e7 53 44 d4 a9 44 49 63 70 08 df 6d ef 85 5c de 59 ee 2c ea d8 e2 23 ce f6 63
2c 86 60 59 d4 35 7a 65 dc b3 ba 5e ef 3b 7f ba 58 46 0b ef 07 b1 fd 1f 86 dc d4 9c
57 b7 57 97 24 49 aa d5 15 e7 5d d6 56 72 f4 d3 53 ad 3c 7a 13 99 d1 68 9e f0 57 59
1a a8 03 0b cf 14 06 06 6c d0 38 81 aa f6 50 3b 7d 6c d8 8e e9 8d f5 94 15 0a 63 85
af 0b 18 28

+Output Parameters:

f5** : de5c85ef6ddf

=====
6.6 Test set 4 [3GPP TS 35.232 V17.0.0 (2022-03)]
=====

IN when computing f5**:
0a 76 87 44 0e ca b6 1e 28 ce 52 e6 a2 27 35 49 1d fb 0e 7c f5 08 6f 44 1f 8e a6 57
b6 6e c1 2b d0 30 2e 31 4b 41 55 54 e8 2b a7 5f 1a 66 c9 86 bd 66 a9 25 54 e5 87 68
00 00 00 00 00 00 00 00 61 cc 6a 4f a1 7d c9 c7 6a 2d 65 50 7a 83 da b8 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 6d b5 df f9 63 30 e1 ae 82 e3 80 fe 5a 86 9e 61
00 80 00 00 00 00
00
00
00 00 00 00
OUT when computing f5**:
93 98 21 4b cc 29 23 f1 75 2b 64 bc 03 b9 e6 14 a3 a9 77 ba 35 f4 a1 13 e7 68 ce a8
01 f3 60 70 e9 0d 4f a3 f6 74 5e 07 7f 0a d2 02 db 67 f6 25 b0 a4 8f f2 2f f6 68 3d
cf f3 dc 79 4a 64 88 62 79 12 d4 ef f1 fe 75 7d ac 07 82 dd 8c 6c ac 00 a9 91 39 f2
69 5e 58 aa 80 86 bf 89 d9 37 ff 24 66 57 e4 11 ba b8 98 82 f5 5c b8 6a 98 c7 4c 59
cc 55 8e 18 fe 6f 08 19 5f e8 06 1f 1a 26 a5 28 9a 47 cb 22 90 ef 5d a2 33 fd 44 79
3c c4 a8 53 87 32 97 75 1d 4e 3d ea 47 4a 67 0c f4 da fd 56 a1 d7 ca 26 25 ad 41 67
57 f7 95 4c 7d e1 96 ce ef e6 8e 25 b9 5f e0 23 92 88 e4 2b 8e 90 9d 78 17 eb 92 4c
af 6c b7 9e

+Output Parameters:
f5** : b8balle45766

=====
6.7 Test set 5 [3GPP TS 35.232 V17.0.0 (2022-03)]
=====

+Intermediate Values:
IN when computing f5**:
62 ff 5e a5 65 61 7d 3e bc 48 cd ae 76 a9 aa f3 e8 23 f2 89 bb 3c aa 47 8a a2 32 15
e4 52 60 3c c9 30 2e 31 4b 41 55 54 ef 8b 49 22 83 08 e3 b1 1f 65 de 8c c6 aa 70 c5
00 00 00 00 00 00 00 00 a5 5a 0e 77 15 9f c2 69 2b aa 26 44 5f 95 44 42 cd c9 89 f7
80 28 c8 89 c1 05 1d 88 56 ca 74 15 15 cb ac 66 2e 1e 02 c6 00 00 00 00 00 00 00 00
00 80 00 00 00 00
00
00 00 00 00
OUT when computing f5**:
2b e1 c9 2a b7 9f 2b cc cb d9 32 97 c7 eb 3c 44 25 20 3f 86 31 b0 04 af 59 74 e7 70
a3 04 b4 bb 74 9f 00 57 42 d6 3e 78 7b ac 48 5f 8e 74 ce fc 99 85 cf f8 ef 2f 06 4e
28 ca 4b 4e 91 74 16 86 ec 0b 81 83 dd 9b 0c 87 8a 9c a0 df a6 b3 bb fb 06 14 b7 92
34 17 d4 10 4f 3e f5 26 76 5d 4d b9 37 a8 eb 86 8e 24 e4 fd d2 dd 4c 81 a4 66 47 b4
b2 e0 7c 5b 35 2b 86 5a 1d 43 73 54 42 77 85 3c b7 15 88 b2 fc b1 7c 1c 89 52 67 eb
1f 65 80 cc 3d b1 7f d6 78 9b 8d 5d 28 00 1c 3c 53 5a 01 d2 4b ba 6a 8f 97 25 93 93
5d ff dd 4a 95 64 06 cc ef da 51 a9 70 93 2d 0f 28 3b 63 62 81 13 05 c6 fd 54 9a 9c
c3 8c 07 e4

+Output Parameters:
f5** : 248e86eba837

=====
6.8 Test set 6 [3GPP TS 35.232 V17.0.0 (2022-03)]
=====

+Intermediate Values:

IN when computing f5** (first call of Keccak core):

f1 7a ec 9c fa 6c 96 49 d1 5c 24 56 7f f4 ec 06 55 b6 9a 17 2a e2 2d c8 d6 fc 62 6c
f2 66 4a b0 e1 30 2e 31 4b 41 55 54 ef 8b 49 22 83 08 e3 b1 1f 65 de 8c c6 aa 70 c5
00 00 00 00 00 00 00 00 a5 5a 0e 77 15 9f c2 69 2b aa 26 44 5f 95 44 42 cd c9 89 f7
80 28 c8 89 c1 05 1d 88 56 ca 74 15 2f dc 58 d4 d9 4a 88 4c 1c b0 3a 8e 63 ac ab 83
75 e8 56 b5 61 ba 3a 06 25 e8 30 ac db 55 73 42 1f 00 00 00 00 00 00 00 00 00 00 00
00
00
00 00 00 00

OUT when computing f5** (first call of Keccak core):

f4 14 08 68 89 23 24 ea 27 07 2b a5 5d 0e 52 fb 0a 69 4e e9 3d 14 0c c8 5f 72 b1 29
2d 05 ca 77 3c c7 cf 7d 72 92 07 94 2d 24 bb 34 b8 68 c9 a9 0b 3e 4d 94 19 e6 f7 5a
ba c7 7d 91 73 cb 97 0a 56 aa ff 96 45 5a 27 25 c4 43 37 c7 11 00 70 71 ca f3 e4 e6
b1 45 cd 9a 98 bf a2 2e 2c a3 de eb 4e c7 a0 14 1e e4 7f e6 6d dd da e7 7e 38 53 e6
68 80 cb ba 5b ba 5f 5e c9 cc 45 f1 46 9b 44 2a f0 90 ea a8 62 e2 2b e5 23 38 16 d6
33 4c 72 96 d9 0c 23 f7 6c fa 99 24 fd 62 1a 58 e4 11 d4 55 f1 db 84 2f a6 7c 80 9b
26 29 51 ec 48 f4 86 0f 68 2f 10 f6 ad 96 57 b4 0b 33 ed 75 db ac 70 fb 4f bd c3 ef
5a e0 f9 d1

IN when computing f5** (second call of Keccak core):

f4 14 08 68 89 23 24 ea 27 07 2b a5 5d 0e 52 fb 0a 69 4e e9 3d 14 0c c8 5f 72 b1 29
2d 05 ca 77 3c c7 cf 7d 72 92 07 94 2d 24 bb 34 b8 68 c9 a9 0b 3e 4d 94 19 e6 f7 5a
ba c7 7d 91 73 cb 97 0a 56 aa ff 96 45 5a 27 25 c4 43 37 c7 11 00 70 71 ca f3 e4 e6
b1 45 cd 9a 98 bf a2 2e 2c a3 de eb 4e c7 a0 14 1e e4 7f e6 6d dd da e7 7e 38 53 e6
68 80 cb ba 5b ba 5f 5e c9 cc 45 f1 46 9b 44 2a f0 90 ea a8 62 e2 2b e5 23 38 16 d6
33 4c 72 96 d9 0c 23 f7 6c fa 99 24 fd 62 1a 58 e4 11 d4 55 f1 db 84 2f a6 7c 80 9b
26 29 51 ec 48 f4 86 0f 68 2f 10 f6 ad 96 57 b4 0b 33 ed 75 db ac 70 fb 4f bd c3 ef
5a e0 f9 d1

OUT when computing f5** (second call of Keccak core):

e9 6f 53 a0 20 5d 59 a9 c6 0b ba 4d 48 cf 00 fe 37 f3 73 ef d7 42 82 94 f1 9c c6 e4
82 e8 1d 9f 5a de 3c f5 b3 d6 f8 80 b9 cd 17 c7 07 8d c2 9b c1 ff a4 e0 fc fe 01 42
54 4e d7 8b 80 1f 51 79 31 9a 53 9d a2 5c 9e 36 70 42 88 43 82 f5 43 b9 6d fc 8b 20
b3 c5 f1 d6 5e 63 4a 14 55 8c bf 40 ae 86 78 d2 14 8e b3 03 60 16 b1 31 70 94 c5 3d
0c 82 b0 0a 39 13 2d f1 cf d5 5d a4 96 97 02 5b 2a 69 a0 00 53 de 6d 85 6b 97 bb 04
32 c1 dd 91 f9 21 2d fb b9 e1 5d 7d cc 87 38 02 5a bd b3 28 f1 40 59 41 17 93 c8 a7
11 dd c4 a0 97 b7 b8 81 30 48 9a 31 73 e0 62 23 79 9e 70 18 e1 18 84 29 f1 8c 45 f4
7a 2c df b4

+Output Parameters:

f5**: 8e14d27886ae

Annex A (informative): Change history

Change history							
Date	Meeting	TDoc	CR	Rev	Cat	Subject/Comment	New version
2025-03	SA#107	SP-250087				Presented for information and approval	1.0.0
2025-03	SA#107					Approved by TSG SA	19.0.0
2025-07	SA#108	SP-250667	0001	-	D	Fixing a reference	19.1.0

History

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V19.1.0	January 2026	Publication